

Paradigm Function Morphology I: The ‘basics’

Andrew Murphy

andrew.murphy@uchicago.edu

Inferential-realizational theories

A word’s association with a particular set of morphosyntactic properties determines a sequence of rule applications defining that word’s inflectional form.

1 Definitions

Paradigm function

A paradigm function is a function which, when applied to the root of a lexeme L paired with a set of morphosyntactic properties appropriate to L, determines the word form occupying the corresponding cell in L’s paradigm.

(1) $\langle \text{BE}, \{1\text{st, singular, present}\} \rangle \rightarrow \text{are}$

$\langle \text{BE}, \{1\text{st, singular, present}\} \rangle$ is a *form/property-set pairing* (FPSP)

Paradigm functions are static well-formedness conditions on cells:

- (2) A cell $\langle W, \sigma \rangle$ in the paradigm of some lexeme L in some language ℓ is well-formed only if ℓ ’s paradigm function relates $\langle W, \sigma \rangle$ to $\langle X, \sigma \rangle$, where X is L’s root.
- (3)
 - a. *Lexeme*: The lexical entry for a word
 - b. *Roots*: A lexeme’s default form, devoid of any inflectional marking
 - c. *Stems*: Any expression to which inflectional exponents may potentially be added
 - d. *Words*: The syntactically free forms occupying the cells of a lexeme’s paradigm
- (4) A set τ of morphosyntactic properties for a lexeme of category C is **WELL-FORMED** in some language ℓ only if τ satisfies the following conditions in ℓ :
 - a. For each property $F: v \in \tau$, $F: v$ is available to lexemes of category C and v is a permissible value for F.
 - b. For any morphosyntactic feature F having v_1, v_2 as permissible values, if $v_1 \neq v_2$ and $F: v_1 \in \tau$, then $F: v_2 \notin \tau$.

Extension

Where σ and τ are well-formed sets of morphosyntactic properties, σ is an **EXTENSION** of τ iff (i) for any atom-valued feature F and any permissible value v for F, if $F: v \in \sigma$; and (ii) for any set-valued feature F and any permissible value ρ for F, if $F: \rho \in \tau$, then $F: \rho' \in \sigma$, where ρ' is an extension of ρ .

(5) $\{\text{TNS: pres, AGR: \{PERS: 1, NUM: pl\}}\}$ is an extension of

- | | |
|--|---|
| a. $\{\text{TNS: pres, AGR: \{PERS: 1, NUM: pl\}}\}$ | f. $\{\text{AGR: \{PERS: 1, NUM: pl\}}\}$ |
| b. $\{\text{TNS: pres, AGR: \{PERS: 1\}}\}$ | g. $\{\text{AGR: \{PERS: 1\}}\}$ |
| c. $\{\text{TNS: pres, AGR: \{NUM: pl\}}\}$ | h. $\{\text{AGR: \{NUM: pl\}}\}$ |
| d. $\{\text{TNS: pres, AGR: \{ \}}\}$ | i. $\{\text{AGR: \{ \}}\}$ |
| e. $\{\text{TNS: pres}\}$ | j. $\{ \}$ |

(6) *Format for paradigm functions:*

$$\text{PF}(\langle X, \sigma \rangle) = \langle Y, \sigma \rangle$$

A concrete example

- (7) $\text{PF}(\langle \text{be}, \{\text{TNS: pres, AGR: \{PERS: 1, NUM: sg\}}\} \rangle) = \langle \text{are}, \{\text{TNS: pres, AGR: \{PERS: 1, NUM: sg\}}\} \rangle$
- (8) $\text{PF}(\langle \text{be}, \sigma \rangle) = \langle \text{are}, \sigma \rangle$, where $\sigma = \{\text{TNS: pres, AGR: \{PERS: 1, NUM: sg\}}\}$

Side remark

What about roots associated with more than one lexeme? e.g. *lie* (‘recline’, ‘speak dishonestly’)

- (9)
 - a. $\text{L-index}(\text{lie}) = \text{LIE}_1$ (‘recline’)
 - b. $\text{L-index}(\text{lie}) = \text{LIE}_2$ (‘speak dishonestly’)

But how do we define what PF is?

2 Realization rules

- (10) *Format for realization rules:*
 $RR_{n, \tau, C}(\langle X, \sigma \rangle) =_{\text{def}} \langle Y, \sigma \rangle$
 a. n is the block to which the rule belongs
 b. τ is the property set that the rule can realize through its application
 c. C indicates the class of lexemes whose paradigms may be defined by this rule
- (11) *Rule argument coherence:*
 $RR_{n, \tau, C}(\langle X, \sigma \rangle)$ is defined iff (a) σ is an extension of τ , (b) $L\text{-index}(X) \in C$, and (c) σ is a well-formed set of morphosyntactic properties of $L\text{-index}(X)$.

How does Stump derive:

- (12) a. Blocking of $-x$ and $-o$ in 3sg aorist
 b. Special forms for $-T, -C?$ / specific 3pl marker for present?
- (13) *Paninian well-formedness condition on inflectional rule blocks:*
 If Q and R are inflectional rules belonging to the same block b , then for any expression X and any complete and well-formed set σ of morphosyntactic properties appropriate to X , either
 a. Q and R are not compatible relative to X and σ , or
 b. Q and R are compatible relative to X and σ and either,
 (i) one is narrower than the other or,
 (ii) there is a third rule in block b which is compatible with Q and R relative to X and σ and is narrower than both Q and R .

How can a paradigm function be defined in terms of RRs?

- (14) $\sigma = \{\text{VFORM: fin, VCE: act, TNS: pres, PRET: no, MOOD: indic, AGR: \{PERS: 1, NUM: pl\}}\}$
- (15) a. $\text{PF}(\langle \textit{krad}, \sigma \rangle) =$
 b. $\text{PF}(\langle \textit{dáva}, \sigma \rangle) =$
- (16) **NARROWEST APPLICABLE RULE:**
 a. $RR_{n, \sigma, C}$ is **NARROWER** than $RR_{n, \tau, C}$ iff σ is an extension of τ and $\sigma \neq \tau$
 b. Where $C = C'$, $RR_{n, \sigma, C}$ is **NARROWER** than $RR_{n, \tau, C}$ iff $C \subseteq C'$.
- (17) *Nar_n notation:*
 Where $RR_{n, \tau, C}$ is the narrowest rule in block n which is applicable to $\langle X, \sigma \rangle$, ' $\text{Nar}_n(\langle X, \sigma \rangle)$ ' represents the result of applying $RR_{n, \tau, C}$ to $\langle X, \sigma \rangle$.

So, how do we define PF?

- (18) *Paradigm function for Bulgarians verbs:*
 Where σ is a complete set of morphosyntactic properties for lexemes of category V , $\text{PF}(\langle X, \sigma \rangle) =_{\text{def}} \text{Nar}_D(\text{Nar}_C(\text{Nar}_B(\text{Nar}_A(\langle X, \sigma \rangle)))$

What if no rule is applicable?

- (19) *Identity Function Default:*
 $RR_{n, \{ \}, \cup}(\langle X, \sigma \rangle) =_{\text{def}} \langle X, \sigma \rangle$

What about the PF giving us $\langle \textit{krádoxme}, \{\text{TNS: aor, PRET: yes, AGR: \{PERS: 1, NUM: pl\}}\} \rangle$?

What about the first singular impf. form?

Rules of referral

How do we get the RoR for 2sg having the same form as 3sg in preterite tenses?

- (20) *Rule of referral (Bulgarian):*
 Where τ is any complete extension of $\{\text{PRET: yes, AGR: \{PERS: 2, NUM: sg\}}\}$, n is any of the rule blocks A to D, and $\sigma' = \sigma / \{\text{AGR: \{PERS: 3\}}\}$,
 $RR_{n, \tau, \nu}(\langle X, \sigma \rangle) =_{\text{def}} \langle Y, \sigma \rangle$, where $\text{Nar}_n(\langle X, \sigma' \rangle) = \langle Y, \sigma' \rangle$.

What does σ/ρ mean?

- (21) a. If $\sigma = \{\text{TNS: impf, PRET: yes, agr: \{PERS: 2, NUM: sg\}}\}$
 and
 b. $\rho = \{\text{AGR: \{PERS: 3\}}\}$
 then
 c. $\sigma/\rho = \{\text{TNS: impf, PRET: yes, agr: \{PERS: 3, NUM: sg\}}\}$

3 Some PFM practice

- (22) *Subject agreement in Nimboran (future tense):*

	singular	dual	plural
1 EXCL	<i>ŋgedúo-d-u</i>	<i>ŋgedúo-k-d-u</i>	<i>ŋgedói⁻¹-d-u</i>
1 INCL	<i>maN-ŋgedúo-d-ám</i>	<i>ŋgedúo-k-d-ám</i>	<i>ŋgedói-k-d-ám</i>
2	<i>ŋgedúo-d-e</i>	<i>ŋgedúo-k-d-e</i>	<i>ŋgedói-k-d-e</i>
3 MASC	<i>ŋgedúo-d-am</i>	<i>ŋgedúo-k-d-am</i>	<i>ŋgedói⁻¹-d-am</i>
3 FEM	<i>ŋgedóu-d-um</i>	<i>ŋgedúo-k-d-um</i>	<i>ŋgedói⁻¹-d-am</i>

What rules/blocks do we need?

References

Stump, Gregory T. (2001). *Inflectional Morphology: A Theory of Paradigm Structure*. Cambridge University Press: Cambridge.